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Dichotomies properties on computational complexity of S -packing coloring problems

Nicolas Gastineau

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Abstract

This work establishes the complexity class of several instances of the S -packing coloring problem: for a graph G , a positive integer k and a non-decreasing list of integers $S = (s_1, \dots, s_k)$, G is S -colorable if its vertices can be partitioned into sets S_i , $i = 1, \dots, k$, where each S_i is an s_i -packing (a set of vertices at pairwise distance greater than s_i). In particular we prove a dichotomy between NP-complete problems and polynomial-time solvable problems for lists of at most four integers.

1 Introduction

We consider only finite undirected connected graphs. For a graph G , an i -packing is a set $X_i \subseteq V(G)$ such that for any distinct pair $u, v \in X_i$, $d_G(u, v) > i$, where $d_G(u, v)$ denotes the distance between u and v . We will use X_i to refer to an i -packing in a graph G . For a nondecreasing sequence of positive integers $S = \{s_i | i > 0\}$, a S - k -(packing)-coloring of G is a partition of $V(G)$ into sets S_1, \dots, S_k , where each S_i is an s_i -packing. For a nondecreasing list of k integers $S' = (s'_1, \dots, s'_k)$, an S' -(packing) coloring of G is an S - k -coloring of G for a sequence S which begins with a list S' . A graph G is S -colorable if there exists an S -coloring of G .

The **S-COL** decision problem consists in determining, for fixed S , if G is S -colorable, for a graph G as input.

By $|S|$, we denote the size of a list S , and by s_i we denote the i th element of a list. Let $S = (s_1, s_2, \dots)$ and $S' = (s'_1, s'_2, \dots)$ be two nondecreasing lists of integers with $|S| = |S'|$. We define an order on the lists by $S \leq S'$ if $s_i \geq s'_i$, for every integer i , $1 \leq i \leq |S|$. Note that if $S \leq S'$, G is S -colorable implies G is S' -colorable.

In this article, for a (s_1, s_2, \dots) -coloring of a graph, we prefer to map vertices to the color multi-set $\{s_1, s_2, \dots\}$ even if two colors can be denoted by the same number. This notation allows the reader to directly see to which type of packing the vertex belongs depending on its color. When needed, we will denote colors of vertices in different i -packings by i_a, i_b, \dots

Let S_d^k be a list only containing k integers d . The problem S_1^k -COL corresponds to the k -coloring problem which is known to be NP-complete for $k \geq 3$. The S -coloring generalizes coloring with distance constraints like the *packing coloring* or the *distance coloring* of a graph. We denote by P-COL, the problem $(1, 2, \dots, k)$ -COL for a graph G and an integer k (with G and k as input). The packing chromatic number [9] of G is the least integer k such that G is $(1, 2, \dots, k)$ -colorable. A series of works [4, 6, 8, 9] considered the packing chromatic number of infinite grids. The d -distance chromatic number [13] of G is the least integer k such that G is S_d^k -colorable. Initially, the concept of S -coloring has been introduced by Goddard et al. [9] and Fiala et al. [7]. The S -coloring problem was considered in other papers [10, 11].

The S -coloring problem, with $|S| = 3$ has been introduced by Goddard et al. [9] in order to determine the complexity of the packing chromatic number when $k = 4$. Moreover, Goddard and Xu [10] have proven that for $|S| = 3$, S -COL is NP-complete if $s_1 = s_2 = 1$ or if $s_1 = 1$ and $s_2 = s_3 = 2$ and polynomial-time solvable otherwise. About the complexity of S -COL, Fiala et al. [7] have proven that P-COL is NP-complete for trees and Argiroffo et al. [1, 2] have proven that P-COL is polynomial-time solvable on some classes of graphs.

In the second section, for a list S of three integers, we determine the family of S -colorable trees. Moreover, we determine dichotomies on cubic graphs, subcubic graphs and bipartite graphs. In the third section, we determine polynomial-time solvable and NP-complete instances of S -COL, for unfixed size of lists. We use these results to determine a dichotomy between NP-complete instances and polynomial-time solvable instances of S -COL for $|S| \leq 4$.

Note that for any nondecreasing list of integers S , we have S -COL in NP.

2 Complexity of S -COL for a list of three integers and several classes of graphs

This section is dedicated to the proofs of two theorems: Theorem 2.1 and Theorem 2.3. In [10], the family of S -colorable graphs, for $|S| = 3$, is described in the case S -COL is polynomial-time solvable. Using the properties of these families of graphs [10], it is easy to determine the S -colorable trees for $|S| = 3$ and $S \neq (1, 2, 2)$. In Theorem 2.1, we determine the $(1, 2, 2)$ -colorable trees by giving a characterization by forbidding subtrees. The following definition gives a construction of this family of forbidden subtrees.

Definition 2.1. Let T_0 and T_1 be the trees from Figure 1. We define a family of trees \mathcal{T} as follows:

- i) $T_0 \in \mathcal{T}, T_1 \in \mathcal{T}$;
- ii) if $T \in \mathcal{T}$, then the tree T' , obtained from T by removing an edge between two vertices u and v of degree 2, by adding four vertices u_1, u_2, u_3 and w and by adding the edges $uu_1, u_1u_2, u_2u_3, u_3v$ and u_2w , is in \mathcal{T} .

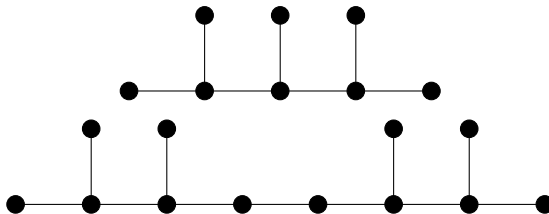


Figure 1: The trees T_0 (on the top) and T_1 (on the bottom).

We can note that the tree T_0 does not contain two adjacent vertices of degree 2. We will prove:

Theorem 2.1. *A tree T is $(1, 2, 2)$ -colorable if and only if it does not contain a tree from \mathcal{T} .*

The following theorem by Goddard and Xu [10] establishes a dichotomy between NP-complete problems and polynomial-time solvable problems for $|S| = 3$.

Theorem 2.2 ([10]). *Let k be a positive integer. The problems $(1, 1, k)$ -COL and $(1, 2, 2)$ -COL are both NP-complete. Except these problems, S -COL is polynomial-time solvable for $|S| = 3$.*

For $(1, 1, k)$ -COL, the authors of [10] provided a reduction from 3-COL with a graph of maximal degree Δ , the produced graph has maximal degree 2Δ . Since 3-COL is polynomial-time solvable for small Δ , the proof can not be easily changed to have a reduction with subcubic graphs. For $(1, 2, 2)$ -COL, the authors provided a reduction from NAE SAT, the produced graph has maximal degree Δ , where Δ is the maximal number of times a variable can appear positively or negatively. This reduction can not be easily changed to have a reduction with subcubic graphs. In Theorem 2.3, we establish similar results for subcubic graphs, cubic graphs and bipartite graphs.

For different instances of S -COL, with $|S| = 3$, Table 1 summarizes the class of complexity of S -COL for different classes of graphs. We recall that every bipartite graph is $(1, 1, k)$ -colorable, for k a positive integer and that every subcubic graph except K_4 is $(1, 1, 1)$ -colorable by Brooks' theorem. We will prove:

Theorem 2.3. *Every instance of S -COL, for $|S| = 3$, is either polynomial-time solvable or NP-complete, for subcubic graphs, cubic graphs and bipartite graphs.*

2.1 Proof of Theorem 2.1

By a 3-vertex we denote a vertex of degree at least 3.

Lemma 2.4. *The trees from \mathcal{T} are not $(1, 2, 2)$ -colorable.*

Class of graphs	(1, 1, 1)-COL	(1, 2, 2)-COL	(1, 1, k)-COL, ($k \geq 2$)
Arbitrary graphs	NP-complete	NP-complete	NP-complete
Subcubic graphs	Polynomial	NP-complete	NP-complete
Cubic graphs	Polynomial	Polynomial	NP-complete
Bipartite graphs	Polynomial	NP-complete	Polynomial
Trees	Polynomial	Polynomial	Polynomial

Table 1: Complexity class of several instances of S -COL.

Proof. First, note that we can not color a 3-vertex by 1. Thus, the tree T_0 , containing three 3-vertices at mutual distance at most 2, is not (1,2,2)-colorable. Second, in the tree T_1 , one of the two vertices of degree 2 should have some color 2. Hence, one vertex among the three 3-vertices at distance at most 2 from this vertex is not colorable. For the other trees, we can note that for every pair of adjacent vertices of degree 2, one of the vertex has some color 2. Moreover, if a vertex is at distance at most 2 from three vertices, then it can not have a color 2. Using these two facts combined together, we can note that if we begin to color two adjacent vertices of degree 3 we cannot extend the (1, 2, 2)-coloring to the other two adjacent vertices of degree 3. □

Proof of Theorem 2.1. By Lemma 2.4, if T contains a tree from \mathcal{T} , then T is not (1,2,2)-colorable. If T does not contain T_0 , then a vertex of degree at least 3 can not be adjacent to two vertices of degree at least 3. Two 3-vertices are *path-connected* if there are only vertices of degree 2 in the path between those two vertices. Let a 3a-vertex (a 3b-vertex, respectively) be a 3-vertex adjacent (not adjacent, respectively) to another 3-vertex.

We construct a (1, 2, 2)-coloring of T as follows: take a 3-vertex u , give an arbitrary color to this vertex and its possible 3-vertex neighbor and extend the coloring to path-connected 3-vertices. Do the same process from these new colored vertices. Let v_1 and v_2 be path-connected 3-vertices at distance i . Suppose, without loss of generality, that v_1 is colored by 2_a . If i is even, then the coloring of v_1 can be extended to v_2 using the pattern $1, 2_b, 1, 2_a, \dots, 1$ for the vertices in the path between v_1 and v_2 . We give the color 2_a (or 2_b , depending on i) to v_2 .

If $i > 3$ is odd then the coloring of v_1 can be extended to v_2 using the pattern $1, 2_b, 2_a, 1, 2_b, 1, 2_a, \dots, 1$ for the vertices in the path between v_1 and v_2 . We give the color 2_a (or 2_b , depending on i) to v_2 .

If $i = 3$, since T does not contain T_1 , then v_1 and v_2 are not both 3a-vertices. First, suppose v_2 is a 3a-vertex. For the two vertices of the path, we give the color 1 to the vertex adjacent to v_1 and the color 2_b to the vertex adjacent to v_2 . Finally, we give the color 2_a to v_2 . Second, suppose v_1 is a 3a-vertex. For the two vertices of the path, we give the color 2_b to the vertex adjacent to v_1 and the color 1 to the vertex adjacent to v_2 . Finally, we give the color 2_a to v_2 . Since G does not contain trees from \mathcal{T} , v_1 is not at distance 2 from another

3b-vertex and v_1 cannot have another neighbor already colored by 2_b . □

2.2 Proof of Theorem 2.3

Definition 2.2. Let S be a nondecreasing list of three integers. Let s_i and s_j be two integers in this list.

An s_i - s_j -transmitter (respectively, an s_i - s_j -antitransmitter) is a graph L such that there exist two vertices u and v called connectors and in every S -coloring of L , u and v have a same color among $\{s_i, s_j\}$ (respectively, different colors among $\{s_i, s_j\}$). Moreover, there should exist an S -coloring of L such that every vertex that have a color c different from the color of u (respectively, of v) must be at distance greater than $\lfloor c/2 \rfloor$ from u (respectively, from v).

An s_i - s_j -clause simulator is a graph H such that there exist three vertices u_1 , u_2 and u_3 called connectors, and in every S -coloring of H , u_1 , u_2 and u_3 have colors among $\{s_i, s_j\}$ and never u_1 , u_2 and u_3 have the same color. Moreover, there should exist S -colorings of H such that every vertex that have a color c different from the color of u_1 (respectively, of u_2 , respectively, of u_3) must be at distance greater than $\lfloor c/2 \rfloor$ from u_1 (respectively, from u_2 , respectively, from u_3), for any possible configuration of colors among $\{s_i, s_j\}$ for the connectors.

Lemma 2.5. Let S be a nondecreasing list of three integers and let s_i and s_j be two integers in this list. If there exist an s_i - s_j -transmitter L , an s_i - s_j -antitransmitter J and an s_i - s_j -clause simulator H , then S -COL is NP-complete.

Proof. Suppose there exist a transmitter L , an antitransmitter J and a clause simulator H . Connecting two vertices u and v by L (respectively, by J) corresponds to adding a copy of L (respectively, of J) and identifying u with a connector of L (respectively, of J) and identifying v with the other connector. The original vertices correspond to the vertices u and v . Let a L - k -cycle be a cycle of order k in which each edge is replaced by L (it means that if uv is an edge, we remove the edge and connect u and v by L).

The proof is by reduction from NAE 3SAT: a literal is either a variable x (a positive occurrence of x) or the negation of a variable (a negative occurrence of x). A clause is a disjunction of literals and the size of a clause is its number of literals. A truth assignment gives values (true or false) for each variable. For a given truth assignment on a set of clauses, if there is a true literal and a false literal in each clause, then it satisfies not equally the set of clauses.

Not-all-equal 3-satisfiability (NAE 3SAT)

Instance: Collection $C = [c_1, c_2, \dots, c_m]$ of clauses on a finite set of variables of size n such that $|c_i| = 3$ for $1 \leq i \leq m$.

Question: Is there a truth assignment for variables that satisfies not equally all clauses in C ?

Let C be a collection of clauses. Let m be the number of clauses and let n be the number of variables. We construct a graph G as follows.

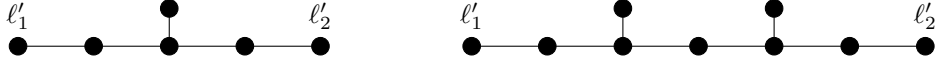


Figure 2: The graph L (on the left) and the graph J (on the right).

For each clause, we associate a clause simulator H with connectors ℓ_1 , ℓ_2 and ℓ_3 representing the three literals of the clause. For each variable x with p_1 positive occurrences of x and p_2 negative occurrences of x , we associate a $L - (p_1 + 1)$ -cycle and a $L - (p_2 + 1)$ -cycle where a connector b_1 of the $L - (p_1 + 1)$ -cycle and a connector b_2 of the $L - (p_2 + 1)$ -cycle are connected by J . For each literal ℓ'_1 which is a positive occurrence of x (respectively, negative occurrence of x), we connect ℓ'_1 by L to a connector of the $L - (p_1 + 1)$ -cycle different from b_1 (respectively, to a connector of the $L - (p_2 + 1)$ -cycle different from b_2), so that every connector of the $L - (p_1 + 1)$ -cycle (respectively, $L - (p_2 + 1)$) is connected to only one occurrence of x .

Suppose there exists a truth assignment of variables that satisfies not equally all clauses in C . To each clause c corresponds a clause simulator H . Given a truth assignment, we give the color s_i to the connectors representing a true literal and s_j to the connectors representing a false literal. If a connector of a transmitter or an antitransmitter is colored, then we give the other possible color to the other connector. Using the existing S -coloring of H , L and J , it gives us an S -coloring of G (the vertices colored have enough distance between them by hypothesis).

Reciprocally, suppose there exists an S -coloring of G . For each H , the connectors of H have colors among $\{s_i, s_j\}$, by hypothesis. By construction, positive occurrences and negative occurrences of a same variable must have different colors. For each positive or negative occurrence of a variable x , we create a truth assignment by giving the value true (respectively, false) to positive occurrences colored by s_i (respectively, s_j) and by giving the value false (respectively, true) to negative occurrences colored by s_i (respectively, s_j). By construction, this truth assignment satisfies not equally all clauses in C . \square

Note that if H , L and J are subcubic graphs, the connectors of H have degree 2 and if the connectors of L and J have degree 1, then the constructed graph is subcubic. Moreover, if H , L and J are bipartite and any pair of connectors of L , J and H are connected by a path of even length, then the constructed graph is bipartite.

Theorem 2.6. *The problem $(1, 2, 2)$ -COL is NP-complete for subcubic bipartite graphs.*

Proof. Let L and J be the subcubic graphs from Figure 2 with the corresponding connectors ℓ'_1 and ℓ'_2 . Remark that if a vertex has degree 3, then it can not be colored by 1, since the coloring can not be extended to a $(1, 2, 2)$ -coloring. Therefore, every connector must be colored by 2_a or by 2_b , since in the final

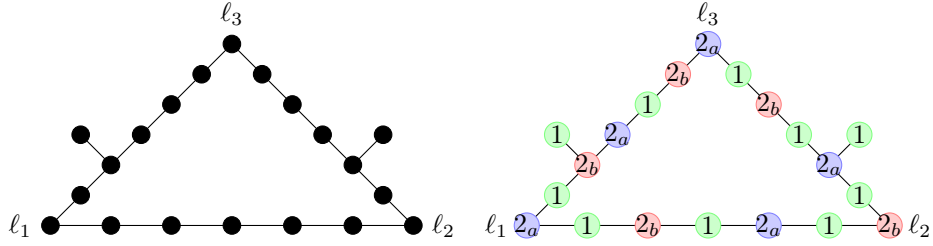


Figure 3: The graph H (on the left) and a $(1, 2, 2)$ -coloring of H (on the right).

constructed graph they will have degree 3. Observe that there are vertices of degree 3 in L and J . Hence, these vertices must also be colored by 2_a or by 2_b .

If the two connectors of L have different colors, then the central vertex of degree 3 can not be colored. Thus, in every $(1, 2, 2)$ -coloring of L , the connectors must have the same color. A $(1, 2, 2)$ -coloring of L exists by giving the color 2_a to the connectors, 2_b to the central vertex of degree 3 and 1 to the remaining vertices.

If the two connectors of J have the same color, then the two central vertices of degree 3 can not be colored. Thus, in every $(1, 2, 2)$ -coloring of J , the connectors must be colored differently. A $(1, 2, 2)$ -coloring of J exists by giving the color 2_a to a connector and to a central vertex of degree 3, 2_b to the other connector and to the other central vertex of degree 3 and 1 to the remaining vertices. To conclude, L is a 2-2-transmitter and J is a 2-2-antitransmitter, respectively.

Let H be the subcubic graph from the left part of Figure 3 with the corresponding connectors ℓ_1 , ℓ_2 and ℓ_3 . The graph H contains two vertices of degree 3 and three connectors. Moreover, these vertices should be colored by 2_a or by 2_b . Suppose ℓ_1 , ℓ_2 and ℓ_3 have the same color 2_a . The two vertices of degree 3 should be colored by 2_b . Observe that the vertices at distance 2 from ℓ_3 could only have the color 1. Thus, the two neighbors of ℓ_3 can not be colored, since they can not be colored by 1 or by 2_a and they can not have both the color 2_b . Therefore, in every $(1, 2, 2)$ -coloring of H the connectors must be colored differently. The right part of Figure 3 gives the existence of an S -coloring with the suited property for ℓ_1 and ℓ_3 with the same color (for the other cases, the $(1, 2, 2)$ -coloring is similar). To conclude, H is a 2-2-clause simulator. \square

Note that no r -regular graph, for $r \geq 3$, is $(1, 2, 2)$ -colorable. Thus, this problem seems to be hard only on graphs containing a large proportion of vertices of degree at most 2.

Theorem 2.7. *Let $k \geq 2$ be a positive integer. The problem $(1, 1, k)$ -COL is NP-complete for subcubic graphs.*

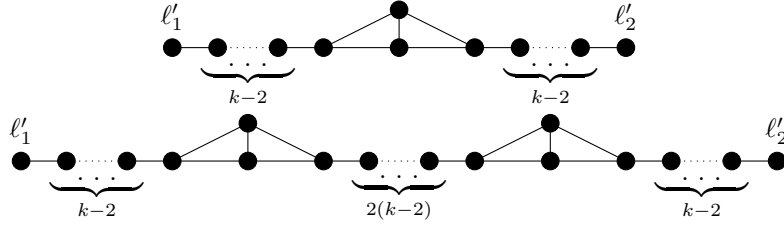


Figure 4: The graph L_k (on the top) and the graph J_k (on the bottom).

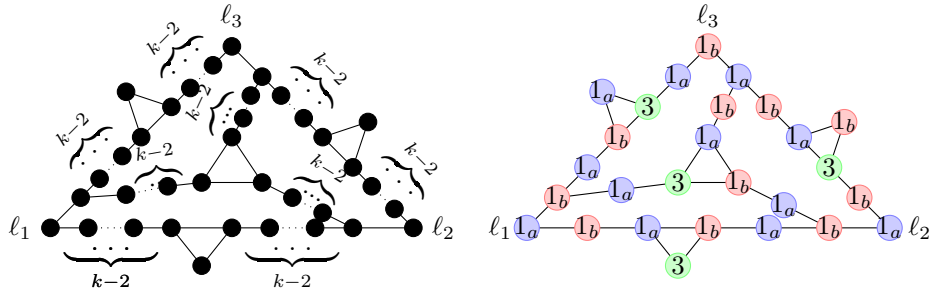


Figure 5: The graph H_k (on the left) and a $(1, 1, 3)$ -coloring of H_3 (on the right).

Proof. Let L_k and J_k be the subcubic graphs from Figure 4 with the corresponding connectors ℓ'_1 and ℓ'_2 . First, observe that there are two triangles in L_k . If no common vertices of the two triangles are colored by k , then the coloring can not be extended to a $(1, 1, k)$ -coloring of L_k . Thus, in every $(1, 1, k)$ -coloring one of these two vertices is colored by k and the remaining vertices could only be colored by 1_a or by 1_b , alternating the colors. Therefore, since the number of the remaining vertices is odd, if one of the connector is colored by a color 1, the other connector could only be colored by the same color 1. In this $(1, 1, k)$ -coloring, the vertices colored by k are at distance k from the connectors.

Observe that there are two pairs of adjacent triangles in J . Thus, for the same reason than for L , for each pair of adjacent triangles, one of the two vertices which is common to both triangles should be colored by k . Therefore, since the number of remaining vertices is even, if one of the connector is colored by a color 1, the other connector could only be colored by the other color 1. To conclude, L_k is a 1-1-transmitter and J_k is a 1-1-antitransmitter.

Let H_k be the subcubic graph from the left part of Figure 5 with the corresponding connectors ℓ_1 , ℓ_2 and ℓ_3 . The graph H_k contains four triangles and in every triangle a vertex should be colored by k . Thus, in every $(1, 1, k)$ -coloring the remaining vertices could only be colored by 1_a or by 1_b alternating the colors. Without loss of generality, suppose ℓ_1 , ℓ_2 and ℓ_3 are colored by the same color

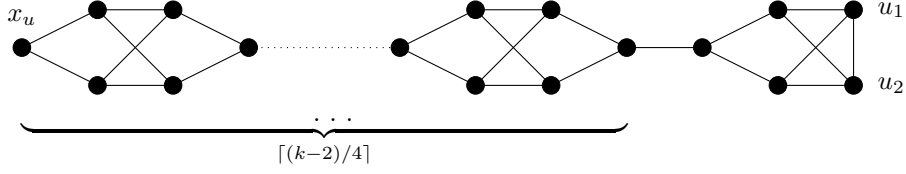


Figure 6: The graph N_k .

1_a . The two neighbors of these vertices should be colored by 1_b and alternating the colors, the central triangle has its neighbors all colored by the same color 1_b or by 1_a (depending on k). Hence, the central triangle can not be colored, the color of the neighbors being not available. The right part of Figure 5 gives the existence of an S -coloring with the suited property for $k = 3$ (for any other integer k , the coloring is similar). To conclude, H_k is a 1-1-clause simulator. \square

Corollary 2.8. *Let $k \geq 2$ be a positive integer. The problem $(1, 1, k)$ -COL is NP-complete for cubic graphs.*

Proof. Let G be a graph and G' be the graph obtained with the reduction from Lemma 2.5 using the 1-1-transmitter, 1-1-antitransmitter and 1-1-clause simulator from Proposition 2.7. For each vertex $u \in V(G')$ of degree 2, we add a graph N_k from Figure 6 and connect u to the vertex x_u of N_k . We obtain a cubic graph in the previous reduction. The graph N_k is clearly $(1, 1, k)$ -colorable, the color k is used to color one vertex among u_1 and u_2 . \square

Proof of Theorem 2.3. Theorem 2.3 is obtained from Theorem 2.6, Theorem 2.7 and Corollary 2.8. \square

3 Dichotomy theorem for $|S| = 4$

This section is dedicated to the proof of our main theorem, a dichotomy property for S -COL when $|S| = 4$.

Theorem 3.1. *Let S be a nondecreasing list of four integers. The problem S -COL is polynomial-time solvable if $S \leq S'$, for $S' = (2, 3, 3, 3)$, $S' = (2, 2, 3, 4)$, $S' = (1, 4, 4, 4)$ or $S' = (1, 2, 5, 6)$, and NP-complete otherwise.*

The section is organized as follows. In the first subsection, we present some properties used to prove that S -COL is polynomial-time solvable. In the second subsection, we present some properties used to prove that S -COL is NP-complete. In the third subsection, we prove Theorem 3.1.

3.1 Polynomial-time solvable instances of S -COL

The following theorem is a result from Bienstock *et al.* [3]. The definitions of pathwidth, treewidth and minors are not given here, the reader can find them in [12, 14]. Note that for a graph G , the treewidth of G is bounded by the pathwidth of G .

Theorem 3.2 ([3]). *For any forest F , every graph without F as minor has a pathwidth bounded by $|V(F)| - 2$.*

We will use the following theorem proven by Fiala *et al.* [7] in order to prove that several instances of S -COL are polynomial-time solvable.

Theorem 3.3 ([7]). *Let S be a nondecreasing list of integers. The problem S -COL is polynomial-time solvable for graphs of bounded treewidth.*

This result has been generalized to the decision problem associated with the S - k -coloring problem (with a graph and k as instances), for a sequence S which contains only integers smaller than a constant [7]. The theorem from Fiala *et al.* [7] uses the result of Courcelle [5] about properties definable in monadic second-order logic.

Let Y_n denotes the graph obtained from taking three copies of P_{n+1} and identifying one end. We start with the following observation:

Claim 3.4. *Let G be a graph and n be a positive integer. If G contains Y_n as minor, then G contains the graph Y_n as subgraph.*

Proposition 3.5. *If there exist a positive integer n such that Y_n is not S -colorable, then S -COL is polynomial-time solvable.*

Proof. We can note that, by Observation 3.4, every graph containing the graph Y_n as minor is not S -colorable. Using Theorem 3.2, every S -colorable graph has treewidth bounded by $3n - 1$. Using Theorem 3.3, S -COL is polynomial-time solvable. \square

We recall the following theorem from Sharp [13] about the complexity of the distance coloring in k colors (S_{d^k} is the list composed of k integers d).

Theorem 3.6 ([13]). *The problem S_{d^k} -COL is NP-complete for $k \geq \lceil 3d/2 \rceil$ and polynomial-time solvable in the other cases.*

Proposition 3.7. *Let S be a nondecreasing list of integers and let $i > 1$ be an integer. If $s_1 = i$, then S -COL is polynomial-time solvable for $|S| < \lceil 3i/2 \rceil$.*

Proof. Using Theorem 3.6, the problem S_{d^k} -COL is polynomial-time solvable for $k < \lceil 3d/2 \rceil$. Every $S_{i|S|}$ -colorable graphs has bounded treewidth [13] and every S -colorable graph is also $S_{i|S|}$ -colorable. Hence, S -colorable graphs have bounded treewidth. Using Theorem 3.3, S -COL is polynomial-time solvable. \square

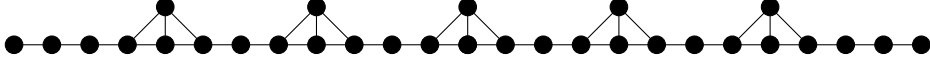


Figure 7: The graph $L_{2,(2,2,2,3)}$.

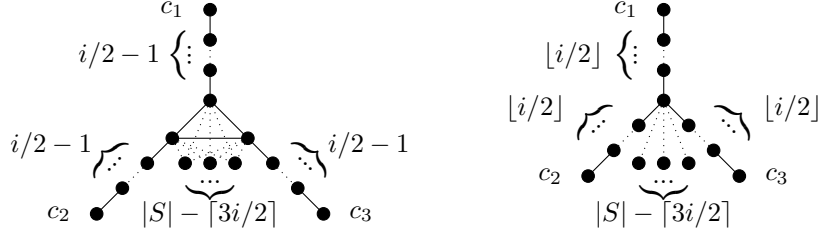


Figure 8: The graphs $R_{i,S}$, for an even i (on the left) and for an odd i (on the right).

3.2 NP-complete instances of S -COL for $|S| \geq 4$

The following proposition is a generalization of the Theorem from Sharp [13].

Theorem 3.8. *The problem S -COL is NP-complete for every list S such that $s_1 = s_{i+1} = i$, with $i > 1$ and $|S| \geq \lceil 3i/2 \rceil$.*

Proof. Let $N_i(S)$ be the number of occurrences of an integer i in the list S and let $n_i^-(S) = N_i(S) - i - 1$. The proof uses a reduction from $N_i(S)$ -COL (as $N_i(S) \geq 3$, $N_i(S)$ -COL is NP-complete).

k -COL

Input: A graph G and an integer k .

Question: Is the graph G k -colorable?

Let G be a graph and $A(S)$ be the following set of integers $\{j > 0 \mid |N_j(S)| \geq 1\}$. We set $m = \prod_{j \in A(S)} (j+1)$. We construct a graph $L_{i,S}$ composed of vertices $u_{j,l}$, for integers j and l , with $-m \leq j \leq m$. Let j, j', l and l' be four integers. Two vertices $u_{j,l}$ and $u_{j',l'}$ in $V(L_{i,S})$ are adjacent if and only if $|j - j'| \leq 1$ (except if $u_{j,l} = u_{j',l'}$). The rank of a vertex is $\text{rank}(u_{j,l}) = j$ and corresponds to its position in $L_{i,S}$. Let j be an integer and let $\text{rank}_{\max}(j) = \max\{l \mid u_{j,l} \in V(L_{i,S})\}$. Let k be an integer, a k -insertion at rank j consists in adding vertices $u_{j,\text{rank}_{\max}(j)+1}, \dots, u_{j,\text{rank}_{\max}(j)+k}$ in the graph $L_{i,S}$.

We construct the graph $L_{i,S}$ as follows: the graph $L_{i,S}$ always contains vertices $u_{-m,0}, \dots, u_{m,0}$. Depending on the list S , we have to do the following insertions:

1. We do an $n_i^-(S)$ -insertion at rank j , for each integer j such that $j \equiv \lfloor (i+1)/2 \rfloor \pmod{i+1}$ and $-m \leq j \leq m$.

2. For every integer $k > i$, we do an $N_k(S)$ -insertion at rank j , for each integer j such that $j \equiv 0 \pmod{k+1}$ and $-m < j < m$.

The last operation consists in adding edges in the graph in order that the adjacency property (two vertices of rank j and j' are adjacent if and only if $|j - j'| \leq 1$) is satisfied. We define *connectors* as specified vertices of degree 1. The connectors of $L_{i,S}$ are the vertices $u_{-m,0}$ and $u_{m,0}$. The graph $L_{2,(2,2,2,3)}$ is illustrated in Figure 7.

Let $R_{i,S}$ be the graph from Figure 8. Note that depending on the parity of i , the graph $R_{i,S}$ is different. The connectors of $R_{i,S}$ are the vertices c_1 , c_2 and c_3 from Figure 8. Connecting a graph G to another graph G' is an operation that consists in identifying a connector of degree 1 of G with a connector of degree 1 of G' . In order to construct the graph $J_{i,S}$, we begin to construct a graph from two copies of $L_{i,S}$ and a copy of $R_{i,S}$ where we connect $R_{i,S}$ to the two copies of $L_{i,S}$. Suppose that u and v are the two connectors of degree 1 of the copies of $L_{i,S}$ in the constructed graph. We finish by deleting one vertex among u and v . If $i = 2$, the following property should be ensured by the copy of $L_{i,S}$: $\text{rank}_{\max}(\text{rank}(w)) = 1$, for w the vertex adjacent to the removed vertex. The constructed graph is the graph $J_{i,S}$. The connectors of this graph are the vertices w and the remaining vertex among u and v .

Let d be a positive integer. The graph G_d is the graph obtained from $R_{i,S}$ and $L_{i,S}$, as follows: we begin with taking d copies L_0, \dots, L_{d-1} of $L_{i,S}$ and d copies R_0, \dots, R_{d-1} of $R_{i,S}$. For every integer k , with $0 \leq k \leq d-1$, we connect R_k to L_k and $L_{k+1 \pmod{d}}$. The obtained graph is the graph G_d and have d connectors, these connectors are the remaining connectors of degree 1 of each copy of $R_{i,S}$.

Finally, we construct the graph G' from the original graph G , as follows: for every vertex $v \in V(G)$, we associate a graph G_v which is a copy of $G_{d(v)}$. We call original vertices, the connectors of a graph G_v , for a vertex v (even if these vertices do not have degree 1 in the constructed graph). For every edge $uv \in E(G)$, we connect G_v and G_u by a graph $J_{i,S}$.

In order to prove that the following construction is a reduction from $N_i(S)$ -COL to S -COL, we prove the following facts:

Fact 1: The graph $R_{i,S}$ is S -colorable and in every S -coloring of $R_{i,S}$, the connectors of $R_{i,S}$ have the same color i .

Proof: The diameter of $R_{i,S}$ is $i+1$ and there are only three vertices which are at pairwise distance at least $i+1$, these vertices are the three connectors of $R_{i,S}$. Since $|V(R_{i,S})| = |S| + 2$, the connectors must have the same color. Moreover, this color must be a color i . There exists an S -coloring of $R_{i,S}$ consisting in giving the same color i to the connectors and giving a different color to each remaining vertex.

Fact 2: For every S -coloring of an induced $L_{i,S}$ in G' and for every color $s_k \neq i$, there is a vertex of rank j colored by s_k , for every integer j such that $j \equiv 0 \pmod{s_k+1}$ and $-m < j < m$.

Proof: There are $\sum_{j=i+1}^{|S|} N_j(S) + 1 = |S| - n_i^-(S)$ vertices of rank 0. Suppose that at least two vertices of rank 0 are colored by some colors i . We have to use

each color i two times in order to color the vertices of rank ℓ , for $-(i+1) < \ell \leq (i+1)$, since there are $|S| + N_i(S)$ vertices of such rank. It is not possible to use each color i two times if two vertices of rank 0 are colored by some color i . Thus, there are vertices of rank 0 of each color $s_k \neq i$.

By induction on j , for $-m \leq j \leq -(i+1)$ and $(i+1) < j \leq m$, we can easily prove that for every color $s_k \neq i$, there is a vertex of rank j colored by s_k , for every integer j such that $j \equiv 0 \pmod{s_k+1}$.

Fact 3: In every S -coloring of $L_{i,S}$ and $J_{i,S}$, the connectors have the same color i in $L_{i,S}$ and two different colors i in $J_{i,S}$.

Proof: Using Fact 2, note that in every S -coloring of $L_{i,S}$ and $J_{i,S}$ there is a vertex of rank j colored by s_k , $s_k \neq i$, for every integer j such that $j \equiv 0 \pmod{s_k+1}$ and $-m < j < m$. Consequently, there is a vertex colored by s_k at distance at most s_k from the connectors of $L_{i,S}$ and $J_{i,S}$. Note also that $-m$ and m have both 0 as remainder modulo $i+1$ in $L_{i,S}$. There is only one vertex colored by i of rank j , for each j such that $j \equiv 0 \pmod{m}$, this color i is the same for each such vertex. Hence, the two connectors of $L_{i,S}$ have the same color i . The argument is similar for $J_{i,S}$ (we have to use Fact 1 also).

Fact 4: There exists an S -coloring of $L_{i,S}$ and $J_{i,S}$.

Proof: First, we color the vertices $u_{0,k}$, for $-m \leq k \leq m$, alternating $i+1$ colors i . Second, for every color $s_k \neq i$, we color the remaining vertices of rank j with color s_k , for every integer j such that $j \equiv 0 \pmod{s_k+1}$ and $-m < j < m$. Finally, we color the remaining vertices of rank j with n_i^- remaining colors i , for $j \equiv \lceil i/2 \rceil \pmod{i+1}$ and $-m \leq j \leq m$. The coloring is similar for $J_{i,S}$.

Fact 5: Let d be a positive integer. In every S -coloring of G_d , the original vertices have the same color i .

Proof: This fact is a consequence of Facts 1 and 3.

Fact 6: For every pair of colors (i_a, i_b) in $N_i(S)$, with $i_a \neq i_b$, there exists an S -coloring of an induced $J_{i,S}$ in G' , with connectors of color i_a and i_b .

Proof: Let L^1 be the copy of $L_{i,S}$ in which no vertex has been removed and let L^2 be the second copy of $L_{i,S}$. We begin by using a list of $i+1$ different colors i to color the vertices $u_{0,k}$, for $-m \leq k \leq m$, in L^1 . We suppose that the color i_a is used to color the initial connectors of the copy of $R_{i,S}$. We give the color i_b to the neighbor of the connector of the copy of $R_{i,S}$ which is a connector of L^2 . We claim that we can extend the coloring to the vertices of the copy of $R_{i,S}$. Using the colorings from Fact 4, we can also extend the coloring to the whole graph (with colors i_a and i_b for the connectors).

Suppose that the graph G is $N_i(S)$ -colorable, with colors $1, \dots, N_i(S)$. For every vertex v of color j , for $1 \leq j \leq N_i(S)$, we use a color s_j in order to color the original vertices of G_v in the graph G' . Using Facts 1, 4 et 6, we can extend the coloring to an S -coloring of the whole graph.

Suppose that G' is S -colorable. Using Fact 5, for each vertex $v \in V(G)$, the original vertices of G_v have the same color i . We conclude that we can construct an $N_i(S)$ -coloring of G using the color j to color vertices with color s_j in G' , for $1 \leq j \leq N_i(S)$. \square

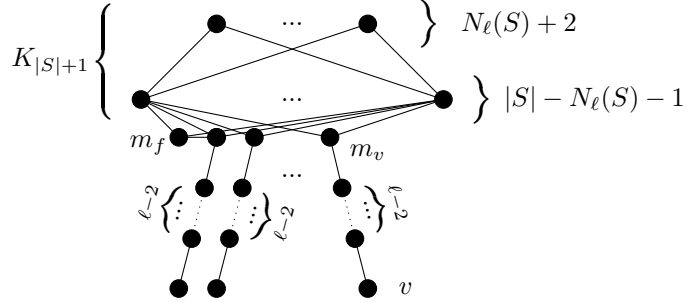


Figure 9: Illustration of the reduction of Proposition 3.10.

Corollary 3.9. *Let k be a positive integer. The problem $(2, 2, 2, k)$ -COL is NP-complete.*

Proposition 3.10. *Let S be a nondecreasing list of integers with $s_2 = 1$ and $|S| \geq 3$. If S -COL is NP-complete, then S' -COL is NP-complete for $|S'| = |S| + 1$, $s'_{|S|+1} \geq s_{|S|}$ and $s'_i = s_i$, for each integer i , with $1 \leq i \leq |S|$.*

Proof. Suppose that S -COL is NP-complete and that $s'_{|S|+1} = \ell$. From a graph G , we construct a new graph G' as follows: for every vertex $v \in V(G)$, we begin with connecting v to the extremity of a path of $\ell - 1$ vertices, with m_v the other extremity of the path (in the case $\ell = 2$, the path has only one vertex and the extremities denote the same vertex). Thus, we add a new vertex m_f that we connect to a vertex m_v , for some v . Afterwards, we add a copy of $K_{|S|+1}$ and we do the join between the set $\{m_v | v \in V(G)\} \cup \{m_f\}$ and a set of $|S| - N_\ell(S) - 1$ vertices in $K_{|S|+1}$. The obtained graph is the graph G' . The original vertices of G' are the vertices that come from the graph G .

Suppose that G is S -colorable. We use an S -coloring of G in order to color the original vertices of G' . Thus, for every vertex v , we alternate the colors $s_1 = 1_a$ and $s_2 = 1_b$ in order to color every path until m_v . Afterwards, we give the color 1_a or 1_b to m_f (depending on the color of the vertex adjacent to m_f). We have $N_\ell(S) + 2$ vertices in $K_{|S|+1}$ non adjacent to m_f . The colors $1_a, 1_b$ and the maximum number of color ℓ in S are given to these vertices. The remaining colors are given to the non colored vertices in $K_{|S|+1}$.

Suppose that G' is S' -colorable. Every vertex of $K_{|S|+1}$ should be colored differently. The two colors given to m_v and m_f should be used in order to color the $N_\ell(S) + 2$ vertices in $K_{|S|+1}$ non adjacent to m_f . Thus, m_v and m_f have some colors 1. At most $N_\ell(S)$ vertices can be colored by ℓ among these $N_\ell(S) + 2$ vertices in $K_{|S|+1}$ non adjacent to m_f . Consequently, a vertex should be colored by a color ℓ among the remaining vertices of $K_{|S|+1}$. This color ℓ can not be used in order to color the other vertices of G' , the other vertices being at distance less than ℓ of a vertex colored by this color ℓ . The coloring of the original vertices of G' is used in order to create an S -coloring of G . \square

Corollary 3.11. *The problem S -COL is NP-complete for every list S such that $s_1 = s_2 = 1$ and $|S| \geq 3$.*

By *subdividing* an edge uv , we mean removing the edge uv , adding a new vertex and connect it to u and v .

Proposition 3.12. *Let S be a nondecreasing list of integers. If S -COL is NP-complete, then S' -COL is NP-complete, for every list S' such that $|S'| = |S| + 1$, $s'_1 = 1$ and $s'_{i+1} = 2s_i$ or $s'_{i+1} = 2s_i + 1$, for every integer i , with $1 \leq i \leq |S|$.*

Proof. Suppose that S -COL is NP-complete. There exists a reduction from S -COL to S' -COL similar with the reduction introduced by Goddard *et al.* [9] from $(1, 1, 2)$ -COL to $(1, 2, 3, 4)$ -COL.

Let G be a connected graph. We construct a new graph G' from G by replacing each edge uv by $|S| + 1$ parallel edges that we subdivide. The vertices u and v are called original vertices.

Suppose that G is S -colorable. Let $X_i \subseteq V(G)$ be an i -packing. Every pair of vertices (u, v) in X_i is such that $d(u, v) > i$. Let $X'_i \subseteq V(G')$ be the set of original vertices in X_i . Observe that every pair of vertices (u, v) in X'_i is such that $d_{G'}(u, v) > 2i + 1$. Thus, X'_i is a $(2i + 1)$ -packing (and consequently a $2i$ -packing). Let X'_1 be the set of non original vertices (created by subdivision). This set is independent and we color the vertices in this set by color 1. We give the color s'_{i+1} to the original vertices of G colored by s_i , for $1 \leq i \leq |S|$, we obtain an S' -coloring of G' .

Suppose that G' is S' -colorable and suppose that there exists an original vertex u colored by 1. Since u has $|S| + 1$ neighbors, two of its neighbors have the same color s'_i . Thus, no original vertex is colored by 1. Using the S' -coloring of the original vertices of G' , we can create an S -coloring of G by giving the color s_{i-1} to the vertices colored by s'_i in G' , for $2 \leq i \leq |S| + 1$. \square

Corollary 3.13. *The problem S -COL is NP-complete for every list such that $|S| = 4$ and $S \geq S'$, with $S' = (1, 3, 5, 5)$.*

3.3 Instances of S -COL for $|S| = 4$

In this section, we prove a dichotomy between NP-complete problems and polynomial-time solvable problems, using the results from the previous section. Let a_0 be the vertex of degree 3 in Y_n and let a_i , b_i and c_i be the eventual vertices at distance i from a_0 in Y_n .

Proposition 3.14. *Let S be a nondecreasing list of integers and let n be an integer. The graph Y_n is not S -colorable for:*

1. $S = (2, 3, 3, 3)$ and $n = 2$;
2. $S = (2, 2, 3, 4)$ and $n = 4$;
3. $S = (1, 4, 4, 4)$ and $n = 3$.

Proof. 1. The vertices a_0, a_1, b_1 and c_1 must have different colors. Thus, one vertex among a_2, b_2 and c_2 can not be colored.
 2. The vertices a_0, a_1, b_1 and c_1 must have different colors. Thus, one vertex among a_4, b_4 and c_4 can not be colored.
 3. Suppose a_0 has color 1. The three vertices a_1, b_1 and c_1 must have different colors. Thus, one vertex among a_2 and a_3 can not be colored. Suppose a_0 has a color 4. Thus, three vertices among a_1, b_1, c_1, a_2, b_2 and c_2 must have a color 4 and one vertex among them can not be colored anymore. \square

Proposition 3.15 ([11]). *The graph P_{16} is not $(1, 3, 5, 8)$ -colorable.*

By adding a leaf on a vertex v , we mean adding a new vertex and connecting it to v . By $\beta(G)$, we denote the order of a *minimum vertex cover* of G .

In the following propositions, we subdivide three times an edge uv . By subdivide three times, we mean replacing an edge by a path of length four with three new vertices. For an edge subdivided three times, the central subdivided vertex is the vertex which is not a neighbor of u or v . We introduce the following operations on a vertex of a graph, the name of these operations come from the color required to have an S -coloring for the introduced vertices:

1. A *2-add* on a vertex u is an operation which consists in adding a vertex w_u and an independent set of vertices I , adding the edge uw_u and adding some edges between $\{u, w_u\}$ and the vertices of I .
2. A *3-add* on a vertex u is an operation which consists in adding a vertex w_u and an independent set of vertices I and adding some edges between $\{u, w_u\}$ and the vertices of I .
3. A *3-5-add* on a vertex u is an operation which consists in doing a 3-add on u and doing several 3-adds or a 2-add on the added vertex w_u .
4. A *3-6-add* on a vertex u is an operation which consists in doing a 3-add on u and doing a 2-add on the added vertex w_u .
5. A *2-5-add* on a vertex u is an operation which consists in doing a 2-add on u with added vertex w_u , doing a 2-add on w_u with added vertex w'_u and doing several 3-adds on the added vertex w'_u .

These operations are illustrated in Figure 10.

Theorem 3.16. *Let $S = (1, 3, k, k')$, with $5 \leq k \leq 7$ and $6 \leq k' \leq 7$. A graph G is S -colorable if and only if one of these three following conditions is true:*

1. $\beta(G) \leq 3$;
2. *the graph G is obtained from a graph G' with vertex cover $C \subseteq \{u, v\}$ of order at least 2 where we can do the following operations:*
 - (a) *if $|C| = 2$ and u and v are adjacent, then we can do several 3-adds on u or v or instead of doing several 3-adds on a vertex among u and v we can do a 2-add on this vertex;*

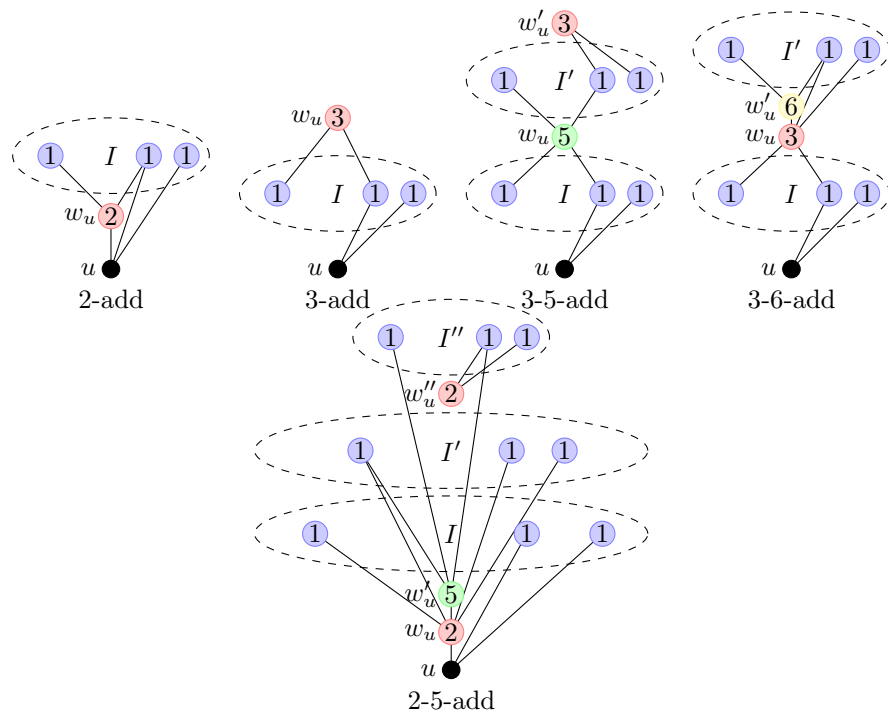


Figure 10: The different used operations in Propositions 3.16 and 3.17 (I , I' and I'' are independent sets).

- (b) otherwise, if $|C| < 1$ or if u and v are not adjacent, then we can do several 3-adds on u or v or instead of doing several 3-adds on u or v we can do a 2-add on these vertices;
3. the graph G is obtained from a bipartite multi-graph $G' = (A, B)$ in which every edge is subdivided three times and where we can do the following operations:
- (a) we can add leaves on the central subdivided vertices and do a 3-add or adding leaves on vertices that come from A or B ;
- (b) if $k \leq 6$, we can do a 3-6-add on vertices that come from A . If $k' \leq 6$, we can do a 3-6-add on the vertices that come from A or B ;
- (c) if $k = 5$, instead of doing a 3-6-add on vertices that come from A , we can do a 3-5-add or several 3-6-adds.

Proof. Let G be a graph verifying one of the three conditions. If the first condition is true, then we color the vertices of the vertex cover by colors 3, k and k' and the remaining vertices by color 1. If the second condition is true, then we color u and v by color k and k' , the other vertices that come from G' by color 1 and the other vertices as in Figure 10. If the third condition is true, then we color the central subdivided vertices by color 3, the vertices that come from A by color k' , the vertices that come from B by color k , the vertices that have been added as in Figure 10 and the remaining vertices by color 1. To conclude, if one of the three conditions is true, then G is S -colorable.

Suppose that G is S -colorable. If no vertex is colored by k or no vertex is colored by k' , then the first condition or the second condition (for a vertex cover of order 1) is true. Hence, we can suppose that there exists at least one vertex of color k and one vertex of color k' . The proof will be organized as follows: we suppose that there exists a vertex u of color k and we consider a vertex v of color k' at minimal distance from u . We will prove that u and v are at distance 4 or that one vertex among u and v can be obtained by doing one of the previous operation on the other vertex. Since the only way to color the vertices in the path between u and v consists in using a color 3 on the central vertex (the vertex at distance 2 from u and v) of this path, the third condition is true about G .

If $d(u, v) = 1$, then the first or the second condition (case (a)) is true. If $d(u, v) = 2$ and if $k > 5$, then the first or the second condition (case (b)) is true. If $d(u, v) = 2$ and $k = 5$, then u is obtained by 3-5-add on v . If $d(u, v) = 3$ and $k = 7$, then the first or the second condition (case (b)) is true. If $d(u, v) = 3$ and $k' \leq 6$ ($k = 6$, respectively), then u (u or v , respectively) is obtained by 3-6-add on v (u or v , respectively). If $d(u, v) > 4$, then the S -coloring can not be extended to the path between u and v . \square

In the following theorem, by an edge uv coming from a multi-graph G , we consider any edge among the multi-edges between u and v .

Theorem 3.17. *Let $S = (1, 2, k, k')$, with $5 \leq k \leq 7$ and $6 \leq k' \leq 7$. A graph G is S -colorable if and only if one of these three following conditions is true:*

1. $\beta(G) \leq 3$;
2. *the graph G is obtained from a graph G' with vertex cover $C \subseteq \{u, v\}$ of order at least 2 where we can do several 3-adds on u or v and doing a 2-add on u or v ;*
3. *the graph G is obtained from a bipartite multi-graph $G' = (A, B)$ where every edge is subdivided three times and where we can do the following operations:*
 - (a) *if $k \leq 6$ and $k' = 7$, we can identify two adjacent subdivided vertices that come from an edge uv where u has only one neighbor in G' and u come from B ;*
 - (b) *if $k = 6$ and $k' = 6$, we can take an independent set of edges from $E(G')$ and identify two adjacent subdivided vertices that come from these edges;*
 - (c) *if $k = 5$ and $k' = 6$, we can take a set of edges from $E(G')$ such that no edge have an extremity in common in B and identify two adjacent subdivided vertices that come from these edges;*
 - (d) *we can add leaves on central subdivided vertices or, if two vertices have been identified, on one of the two subdivided vertices. We can never add leaves on two vertices at distance 2. We can do 3-adds, a 2-add and add leaves on vertices that come from A or B . We can not do a 2-add on a vertex which is adjacent to a vertex coming from an identification;*
 - (e) *if $k = 5$, instead of doing a 2-add on vertices that come from A , we can do a 2-5-add or a 3-5-add on vertices that come from A . We can not do a 2-5-add on vertices at distance at most 2 from a vertex coming from an identification.*

Proof. Let G be a graph verifying one of the three conditions. If the first condition is true, then we color the vertices of the vertex cover by colors 2, k and k' and the remaining vertices by color 1. If the second condition is true, then we color u and v by color k and k' , the other vertices that come from G' by color 1 and the other vertices as in Figure 10. If the third condition is true, then we color the central subdivided vertices by color 2, the vertices that come from A by color k' (if two subdivided vertices have been identified, we color the possible vertex on which leaves have been added by color 2), the vertices that come from B by color k , the vertices that have been added as in Figure 10 and the remaining vertices by color 1. To conclude, if one of the three conditions is true, then G is S -colorable.

Suppose that G is S -colorable. If no vertex is colored by k or no vertex is colored by k' , then the first condition or the second condition is true. Hence,

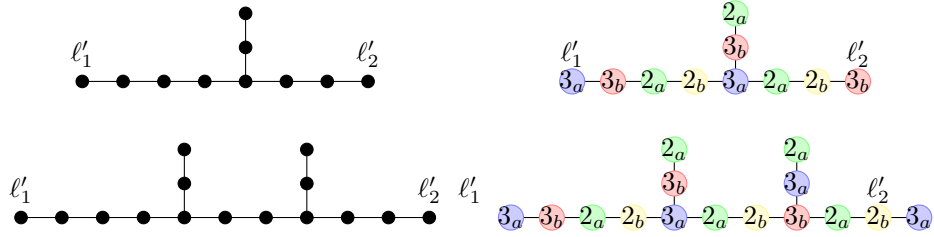


Figure 11: The graphs L' (at the bottom) and J' (at the top) and a $(2, 2, 3, 3)$ -coloring of L' and J' (on the right).

we can suppose that there exists at least one vertex of color k and one vertex of color k' . The proof will be organized as follows: we suppose that there exists a vertex u of color k and we consider a vertex v of color k' at minimal distance from u . We will prove that u and v are at distance between 3 and 4 or prove that one vertex among u and v can be obtained by doing one of the previous operation on the other vertex. Since the only way to color the vertices in the path between u and v consists in using a color 2 on one of the vertex of this path (the possible vertex of degree at least 3 or the central vertex if no vertices have been identified), the third condition is true about G .

If $d(u, v) = 1$, then the first or the second condition is true. If $d(u, v) = 2$ and if $k > 5$, then the first or the second condition is true. If $d(u, v) = 2$ and $k = 5$, then u is obtained by 2-5-add or by 3-5-add on v . If $d(u, v) > 4$, then the S -coloring can not be extended to the path between u and v . \square

In the following proof, we use the construction from Lemma 2.5. The existence properties of an S -coloring with the properties of Lemma 2.5 will not be verified anymore. However, using the different $(2, 2, 3, 3)$ -colorings from the proof, the reader can be convinced that there exists an S -coloring of the whole constructed graph.

Theorem 3.18. *The problem $(2, 2, 3, 3)$ -COL is NP-complete.*

Proof. Let L' and J' be subcubic graphs from Figure 11 with connectors the vertices ℓ'_1 and ℓ'_2 . Note that a vertex of degree at least 3 can only be colored by colors 3_a or 3_b . In the constructed graph (Lemma 2.5), a vertex of degree at least 3 will be in an induced Y_2 and there does not exist an S -coloring of Y_2 with the central vertex colored by colors 2_a or 2_b . For the same reason, every connectors should be colored by colors 3_a or 3_b (these vertices will be in an induced Y_2 in the constructed graphs). If two vertices u and v have degree at least 3 or are connectors and $d(u, v) = 4$, then u and v can only have the same color 3, since the vertices in the path between u and v can not be colored anymore if u and v have different colors.

If two connectors of L' have the same color, then the vertices in the path of length 4 between the central vertex of degree 3 and a connector can not be

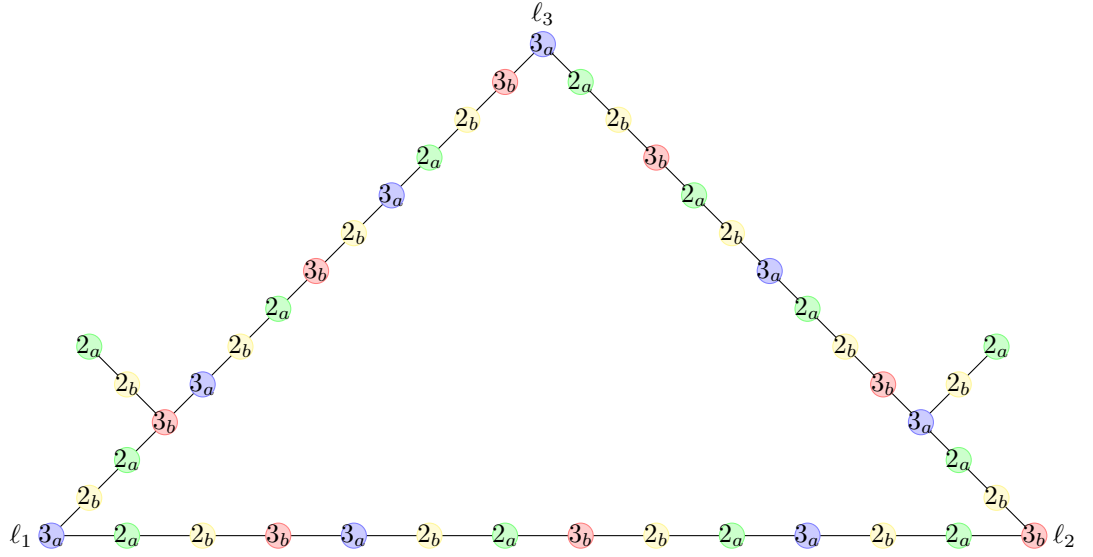


Figure 12: The graph H' and a $(2, 2, 3, 3)$ -coloring of it.

colored. Thus, in every $(2, 2, 3, 3)$ -coloring of L' , the connectors should have the same color 3. A $(2, 2, 3, 3)$ -coloring of L' is illustrated in Figure 11.

If two connectors of J' have different colors 3, then the vertices in the path of length 4 between one of the central vertex of degree 3 and a connector can not be colored. Thus, in every $(2, 2, 3, 3)$ -coloring of J' , the connectors can not have the same color 3. A $(2, 2, 3, 3)$ -coloration of J' is illustrated in Figure 11. Hence, L' is a 3-3-transmitter and J' is a 3-3-antitransmitter.

Let H' be the subcubic graph from Figure 12 with connectors the vertices ℓ_1, ℓ_2 and ℓ_3 . The graph H' contains three connectors and two vertices of degree 3. These vertices should be all colored by colors 3_a or 3_b . Suppose that ℓ_1, ℓ_2 and ℓ_3 have the same color 3_a . Let u be one of the vertex at distance 3 from ℓ_1 . Suppose that u is colored by color 3_b . The coloring can not be extended to the vertices in the path between u and the vertex of degree 3 at minimum distance of u , since no vertices in this path can be colored by 3_b and P_6 is not $(2, 2, 3)$ -colorable. Thus, the vertices at distance 3 from ℓ_1 can not be colored by 3_b . One vertex among u and the vertices in the shortest path between u and ℓ_1 must be colored by 3_b . There does not exist a $(2, 2, 3, 3)$ -coloring of Y_2 for which a vertex at distance 2 from a vertex of degree 3 is colored by a color 3. Thus, the only way to color H is to color the neighbors of ℓ_3 by color 3_b . Hence, we have a contradiction (the distance between these two vertices of color 3_b is only 2). Consequently, in every $(2, 2, 3, 3)$ -coloring of H' , the connectors can not have a same color.

Figure 12 illustrates a $(2, 2, 3, 3)$ -coloring for ℓ_1 and ℓ_3 with the same color

(in the other cases, the $(2, 2, 3, 3)$ -colorings are similar). To conclude, H' is a 3-3-clause simulator. \square

Proof of Theorem 3.1. By Propositions 3.7, 3.14, 3.15, 3.16 and 3.17, we obtain that S -COL is polynomial-time solvable in the given cases. By Propositions 3.9, 3.11 and 3.13, we obtain that S -COL is NP-complete in the other cases. \square

4 Conclusion

We have proven several dichotomy theorems about S -COL. For several instances of S -COL ($(2, 3, 3, 3)$ -COL, for example) the family of S -colorable graphs is included in some known families of graphs, since every S -colorable graph has bounded pathwidth. For other instances of S -COL ($(2, 2, 3, 3)$ -COL, for example), S -COL is NP-complete, by giving a reduction. For a last kind of instances of S -COL ($(1, 2, 7, 7)$ -COL, for example), we determine a construction of the S -colorable graphs. The structure of these S -colorable graphs is not so far from the structure of a bipartite graph.

We finish by proposing the following conjecture about S -COL.

Claim 4.1. *There exists a dichotomy between NP-complete problems and polynomial-time solvable problems among the instances of S -COL.*

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